# Week 6: The Quantum Fourier Transform

COMS 4281 (Fall 2024)

- 1. Practice worksheet and quiz #3 will be out tonight.
- 2. Midterm on October 21. More details soon.

- Quantum circuits
- Universal gate sets, Solovay-Kitaev Theorem
- Deutsch and Simons algorithms

Quantum algorithm that implements the Discrete Fourier Transform (DFT) on exponentially large dimension.

It is the heart of many powerful quantum algorithms such as

- Shor's factoring algorithm
- Phase estimation algorithm
- Algorithms for solving hidden subgroup problem

# **Discrete Fourier Transform**

A method to uncover **hidden**, **periodic structure** in vectors. Used everywhere in engineering, science, and mathematics.

An example of a vector that represents a *noisy signal* whose characteristics we'd like to analyze.



The entries of the vector correspond to equally spaced time points and the value of the entry corresponds to the signal amplitude at that time. The **Discrete Fourier Transform (DFT)** is a method to express every vector (i.e. every signal) as a linear combination of simple periodic vectors (i.e. complex sinusoidal signals).

Visually:



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**Mathematically**: every vector  $|\psi\rangle \in \mathbb{C}^N$  can be written as

$$|\psi\rangle = \sum_{j=0}^{N-1} \hat{\psi}_j |f_j\rangle$$

where  $\{|f_0\rangle, |f_1\rangle, \dots, |f_{N-1}\rangle\}$  is the  $\mathbb{Z}_N$ -Fourier basis.

The Fourier basis vectors are

$$|f_j\rangle = rac{1}{\sqrt{N}}\sum_{k=0}^{N-1}\omega_N^{jk} |k
angle$$

for  $j = 0, 1, \ldots, N-1$  where

$$\omega_{N} = \exp\left(2\pi i/N\right)$$

are the N'th roots of unity.



Example: N = 8, j = 0.

$$|f_0\rangle = \frac{1}{\sqrt{N}}\sum_{k=0}^{N-1}|k\rangle = \frac{1}{\sqrt{N}}\Big(1,1,\ldots,1\Big)^{\top}$$



Example: N = 8, j = 1.

$$|f_1\rangle = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} \omega_N^k |k\rangle = \frac{1}{\sqrt{N}} \left(1, \omega_8, \omega_8^2, \dots, \omega_8^7\right)^\top$$

Visually, the real component of  $\omega_N^{jk}$ :



Example: N = 8, j = 5.

$$|f_5\rangle = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} \omega_N^{5k} |k\rangle = \frac{1}{\sqrt{N}} \Big( 1, \omega_8^5, \omega_8^2, \omega_8^7, \dots \Big)^\top$$

Visually, the real component of  $\omega_N^{jk}$ :



- Let  $|\psi\rangle = \sum_{j=0}^{N-1} \psi_j |j\rangle$  be a signal represented in the standard basis.
- In the Fourier basis,  $|\psi\rangle$  can be written as  $|\psi\rangle = \sum_{j=0}^{N-1} \hat{\psi}_j |f_j\rangle$  for some Fourier coefficients  $\hat{\psi}_j$ .

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- The magnitude  $|\hat{\psi}_j|^2$  of the *j*'th Fourier coefficient quantifies the contribution of  $|f_j\rangle$  to  $|\psi\rangle$ .
- What's the relationship between the amplitudes  $\{\psi_0, \psi_1, \dots, \psi_{N-1}\}$  and the Fourier coefficients  $\{\hat{\psi}_0, \hat{\psi}_1, \dots, \hat{\psi}_{N-1}\}$ ?

They are related by a *unitary* transformation  $F_N$  (which is the DFT). It maps  $|j\rangle \mapsto |f_j\rangle$ . The *inverse* DFT is  $F_N^{\dagger}$ , which maps  $|f_j\rangle$  to  $|j\rangle$ .

Applying  ${\cal F}_{\cal N}^\dagger$  to  $|\psi
angle$  yields the vector

$$F_{N}^{\dagger} |\psi\rangle = |\hat{\psi}\rangle = \sum_{j=0}^{N-1} \hat{\psi}_{j} |j\rangle.$$

The Fourier coefficients of  $|\psi\rangle$  have been turned into *amplitudes* in the standard basis of  $|\hat{\psi}\rangle$ .

# **Quantum Fourier Transform**

Let  $N = 2^n$ . The **Quantum Fourier Transform (QFT)** is a quantum algorithm for implementing the *n*-qubit unitary  $F_N$  while taking only poly(n) time.

QFT maps quantum states  $|\psi\rangle$  to their Fourier transform state  $|\hat{\psi}\rangle=\sum_{j}\hat{\psi}_{j}\,|j\rangle.$ 

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QFT maps quantum states  $|\psi\rangle$  to their Fourier transform state  $|\hat{\psi}\rangle = \sum_{j} \hat{\psi}_{j} |j\rangle$ .

Here, we associate  $|j\rangle$  with the *n*-qubit state  $|j_1j_2\cdots j_n\rangle$ , the binary representation of *j*.

The best classical algorithm for computing DFT of a N-dimensional vector is called the **Fast Fourier Transform** (**FFT**), which takes  $O(N \log N)$  time.

Does this constitute an exponential speedup?

- Not exactly. The Fast Fourier Transform gets an input that is a **classical** N-dimensional vector  $|\psi\rangle$ , and outputs another N-dimensional vector  $|\hat{\psi}\rangle$ .
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On the other hand, QFT gets an input vector in **quantum form**, and produces an output vector in quantum form.

The Fourier coefficients  $\{\hat{\psi}_j\}_j$  are not readily accessible: measuring  $|\hat{\psi}\rangle$  produces basis vector  $|j\rangle$  with probability  $|\hat{\psi}_j|^2$ , but afterwards all other information about the Fourier coefficients is lost.

### Example: N = 2 (single-qubit unitary)

$$F_2 = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

.

We've seen this before...

Example: N = 4 (two qubits). Reordering the rows/columns according to  $|00\rangle$ ,  $|10\rangle$ ,  $|01\rangle$ ,  $|11\rangle$ .

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$$F_4 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1\\ 1 & -1 & i & -i\\ 1 & 1 & -1 & -1\\ 1 & -1 & -i & i \end{pmatrix}$$

Can you spot recursive structure?

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$$F_4 = \frac{1}{\sqrt{4}} \begin{pmatrix} 1 & 1 & 1 & 1\\ 1 & -1 & i & -i\\ 1 & 1 & -1 & -1\\ 1 & -1 & -i & i \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} F_2 & A_2 F_2\\ F_2 & -A_2 F_2 \end{pmatrix}$$

where

$$A_2 = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}$$

For general  $N = 2^n$ , if we order the columns where all the "even" columns are on the left, and all the "odd" columns are on the right, then

$$F_{N} = \frac{1}{\sqrt{2}} \begin{pmatrix} F_{\frac{N}{2}} & A_{\frac{N}{2}}F_{\frac{N}{2}} \\ F_{\frac{N}{2}} & -A_{\frac{N}{2}}F_{\frac{N}{2}} \end{pmatrix}$$

where

The recursive formula for  $F_N$  inspires a recursive construction for the QFT circuit. Assume we already have circuits for the unitaries  $F_{N/2}$  and  $A_{N/2}$ . Then the the circuit for  $F_N$  looks like



#### How is the circuit for $A_{N/2}$ implemented?

Note that the unitary acts on n-1 qubits, and for all  $y \in \{0,1\}^{n-1}$  the unitary maps

$$|y_1,\ldots,y_{n-1}\rangle\mapsto\omega_N^{\operatorname{toint}(y)}|y_1,\ldots,y_{n-1}\rangle$$

where

$$toint(y) = y_1 2^{n-2} + y_2 2^{n-3} + \dots + y_{n-1}$$

## Expanding and regrouping we get

$$\begin{aligned} |y_1, \dots, y_{n-1}\rangle &\mapsto \omega_N^{2^{n-2} \cdot y_1} \omega_N^{2^{n-3} \cdot y_2} \cdots \omega_N^{y_{n-1}} |y_1, \dots, y_{n-1}\rangle \\ &= \left(\omega_N^{2^{n-2} \cdot y_1} |y_1\rangle\right) \left(\omega_N^{2^{n-3} \cdot y_2} |y_2\rangle\right) \cdots \left(\omega_N^{y_{n-1}} |y_{n-1}\rangle\right) \end{aligned}$$

#### It is just a tensor product unitary!

$$A_{N/2} = P(1/4) \otimes P(1/8) \otimes \cdots \otimes P(1/2^n)$$

where

$${\it P}(arphi) = egin{pmatrix} 1 & 0 \ 0 & e^{2\pi i arphi} \end{pmatrix}$$

**Complexity analysis**: Let T(n) denote the number of gates used to construct  $F_N$ . It satisfies T(n) = T(n-1) + O(n). Unrolling the recursion, we get  $T(n) = O(n^2)$ .



Why does it work?

It's a couple pages of algebra... take a look at the scribe notes, or in Nielsen/Chuang Chapter 5 if you're interested!

Phase Estimation, the RSA Cryptosystem, and Shor's Algorithm.